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An Efficient Compressive Sensing Data Gathering Using Modified Ant Colony and Diffusion Wavelets in WSN

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Abstract— Compressive sensing (CS) depend on data gathering is a promising method to reduce energy consumption in wireless sensor networks (WSNs). The existing CS-based data-gathering approaches require a large number of sensor nodes to participate in each CS measurement task, resulting in high energy consumption, and do not guarantee load balance. The propose a sparser analysis that depends on modified diffusion wavelets, which exploit sensor readings' spatial correlation in WSNs. In particular, a novel data-gathering scheme with combine routing and CS is present. A modified ant colony algorithm-based diffusion wavelets (ACBDW), where next hop node selection takes a node's residual energy and path length into consideration simultaneously. Moreover, in order to quickness up the coverage rate and avoid the local optimal of the algorithm, an improved pheromone impact factor is put forward. The diffusion wavelets based on sensor nodes' degree and different nodes' distance considering the above factors are proposed. To further reduce the transport costs in WSNs, a sparse measurement matrix is utilized and modified ant colony routing are jointly applied to mitigate energy consumption and balance the network load, especially lowering the transmission costs for those nodes nearest the sink node. The experimental result show data-gathering approaches, this proposed algorithm not only minimizes the energy consumption of the network, but prolongs the network lifetime.

Keywords— Compressive Sensing, Data Gathering, Modified Diffusion Wavelets, Ant Colony Algorithm Spatial Property

I. INTRODUCTION

Wireless sensor networks (WSNs) have received significant attention due to their versatility and have been deployed widely in applications such as military surveillance, monitoring of environment, traffic and critical infrastructures, among others. Increasing the lifetime of a wireless sensor network depends directly on minimizing the energy consumption at sensor nodes. In a WSN, most of the power is consumed in data transmission and forwarding when compared to data sensing and computation (processing). Data collection from impassable terrain and then transmit the information to the sink is a fundamental task in periodic sensor

networks. However, in order to keep these networks operating for long time, adaptive sampling approach to periodic data collection constitutes a fundamental mechanism for energy optimization. The key idea behind this path is to allow each sensor node to adapt its sampling rates to the physical changing dynamics. In this way, over-sampling can be minimized and power efficiency of the overall network system can be further improved. Due to the large amounts of data generated by wireless sensor networks (WSN), collection of sensing data to be forwarded from the sensor nodes to a central base station constitute a major process for WSN[1]. This process is known as Data Collection. Unfortunately, the transmission of large quantities of data is a threat on the lifetime of the sensor network due to the limited energy resources of the sensor nodes. On the other hand, most of the applications require all the data generated and doesn't tolerate any loss of detail as to reach the accuracy required. The "periodic sampling" data collection model is characterized by the acquisition of sensor data from a number of remote sensor nodes then pushing them to the sink on a periodic basis. These periodic models are used for applications where certain conditions or processes need to be monitored constantly, such as the temperature or pressure, etc. Commonly, the data compression techniques can be divided into two different types; lossless and lossy compression. Lossless compression approach, as the name implies involve no loss of information. In other words, the original data can be recovered exactly from the compressed data. This can be obtained by employing the statistical redundancy to represent the sender's data more with fewer errors. In contrast, lossy compression techniques involve some loss of information and data that have been compressed using lossy technique generally cannot be recovered or reconstructed exactly. From other view point, the data compression techniques in WSNs can be classified into five categories: (1) the string-based compression techniques treat sensing data as a chain of characters and then approve the text data compression schemes to compress them. (2) The image-based compression techniques hierarchically organize WSNs and then employ the idea from the image compression solutions to handle sensing data. (3) The distributed source coding approach extend the Slepian-Wolf theorem to encode multiple correlated data streams independently at sensor nodes and then jointly decode them at the sink. (4) The compressed sensing techniques adopt a small number of non-adaptive and randomized linear projection samples to compress sensing data. (5) The data aggregation approach select a subset of sensor nodes in the network to be responsible for fusing the sensing data from other sensor nodes to reduce the amount of data transmissions. Advances in computing and communication technologies have led to intensive research effort on wireless sensor networks (WSNs). WSNs have found extensive applications in urban traffic monitoring and environmental surveillance. Typically, a WSN consists of a number of sensor nodes, which are randomly distributed in the field under surveillance, and a sink node. Generally, sensor nodes are required to collect data periodically and transmit them to the sink through multi-hop routing, and then the information aggregation and extraction tasks are performed at the sink [2]. Considering that sensor nodes usually have limited energy supply and that replacing or recharging the batteries of sensor nodes is difficult in practical WSN deployments, a primary objective of data gathering in WSNs is to obtain an accurate approximation of the signal field with as little energy expenditure as possible. Rest of the paper is organized as follows, section I contain overview of wireless sensor network and data collection model. Section II contain review of exiting compressive data gathering algorithms, Section III contain proposes system and module implementations, Section IV contain result and discussion, performance analysis, Section V concludes.

II. RELATED WORK

For the problems of random selection and unbalanced position of projection nodes, this paper proposes a compressed data gathering method based on even projection. For the WSN with uniformly distributed nodes, a location-based even clustering method is proposed. The clustering is implemented with the same size of the grids, which ensures the positional balance of the projected nodes. For the WSN with uneven distributed nodes, a node density-based even clustering method is proposed. The DEC method, taking into account the factors of location and density, reduces the energy consumption at isolated points, equalizes the energy, and extends the network lifetime. Moreover, the analysis and simulation of the relevant parameters affecting the network energy consumption were analyzed. Compared with the random projection node method and the random walk method, the proposed method performs well and the network lifetime is significantly extended. In the next step, we will consider the application of artificial intelligence to further optimize the routing topology of the network and make more in-depth research on signal reconstruction to obtain better compressed data collection results. The data gathering method by combining the Compressive Data Gathering (CDG) presented with sparse random projection presented to reduce further the overall number of transmissions and most importantly to distribute the energy consumption load more evenly throughout the network to increase the lifetime of wireless sensor network. Our method (the Minimum Spanning Tree Projection (MSTP)), same as selects different nodes at random to do projection. Where in each projection node after collecting the native data from set of nodes sends the projected data to the sink. But MSTP unlike uses CDG for each projection node to collect and gathers one weighted sum by constructing independent forwarding tree which ensures fewer transmissions. The data gathering is able to reduce global scale communication cost without introducing intensive computation or complicated transmission control. The load balancing essential is capable of extending the lifetime of the entire sensor network as well as individual sensors.

Understanding of recovering a given sparse signal with sparse random matrices in the presence of channel fading. More specifically, we provide lower bounds on the number of measurements that should be collected by the fusion center in order to achieve non uniform recovery guarantees with l_1 norm minimization-based recovery with independent (not necessarily identical) fading channels. With sparse random projections, the nodes transmit their observations with a certain probability[3]. We further discuss how to design probabilities of transmissions by each node (equivalently the sparsity parameter of the random projection matrix) based on the channel fading statistics so that the number of measurements required for signal recovery at the fusion center is minimized. consider a similar problem of WCS. A distributed compressive sensing arrangement for WSNs in order to reduce computational complexity and the communication cost. It considers an $M \times N$ sparse

random matrix with entries that have a probability g of being nonzero, so that on average there are ng non-zeros per row. The resulted data similarity error rate is comparable to that of the optimal k -term approximation if the energy of the signal is not concentrated in a few elements. Somehow, the sparsity component g of random projections impacts the accuracy of signal reconstructions. Usually, the sparsity factor g is statistically determined according to the amount of harvested energy and is homogeneous for all sensors only considered AWGN channels. Compressive Sensing-based clustering algorithm, a clustered WSN only needs to send M measurements from its clusters to the BS. All raw reading data from N sensors will be recovered based on those measurements at the BS. The algorithm helps to reduce a significant energy consumption to transmit data from the network to the BS. Furthermore, we formulate the total power consumption for clustered WSNs that apply the algorithm[4]. We

analyze the total power consumption of the network versus number of clusters. Both common positions of the BS are considered: the BS at the center and outside the sensing area. Based on that, we can obtain the optimal number of clusters that provides the minimum power consumption for our networks. The energy-efficient data collection in wireless sensor networks (WSNs) that is based on an integration of the clustering and compressive sensing (CS). It is well known that natural signals have spatial correlation and therefore the sensor readings in a WSN are sparse in a proper basis such as DCT or wavelet. Cluster-based routing strategy has several advantages such as conserving communication bandwidth, stabilizing the network topology and reducing the rate of energy consumption. The major factors influencing the energy consumption of the clustering scheme are the number of clusters and the distribution of cluster heads.[5] Optimizing the number of clusters in WSN has been addressed by many researchers. But these methods are based on conventional in-network compression and/or non-compression data gathering. Weighted Compressive Data gathering (WCDA)”, which benefits from the advantage of the sparse random measurement matrix to reduce the energy consumption. The novelty of the WCDA algorithm lies in the power control capacity in sensor nodes to form energy efficient routing trees with focus on the load-balancing issue. In the second part, we present another new data aggregation method namely “Cluster-based Weighted Compressive Data Aggregation (CWCDA)” to make a significant reduction in the energy consumption in our WSN model. The main idea after this algorithm is to apply the WCDA algorithm to each cluster in order to reduce significantly the number of involved sensor nodes during each CS measurement. Hierarchical Data Aggregation method using Compressive Sensing (HDACS) is presented, which combines a hierarchical network configuration with CS[6]. Our key idea is to set multiple compression thresholds modified based on cluster sizes at different levels of the data aggregation tree to optimize the amount of data transmitted.

III. PROPOSED METHODOLOGY

Wireless Sensor Networks (WSNs) generally consist of a large number of sensor nodes and a sink node deployed in the detected environment to monitor various physical characteristics of the real world, such as temperature, voltage, wind direction, and so on. Furthermore, WSNs should have a long enough lifetime to successfully fulfill the monitoring task. However, sensor nodes are limited in terms of computational ability, communication bandwidth, and energy availability.[7] The intuition behind CDG is that higher efficiency can be achieved if correlated sensor readings are transmitted jointly rather than separately. Showing how sensor readings are incorporate while being relayed along a chain-type topology to the sink. In practice, sensors usually increase in a two-dimensional area, and the ensemble of routing paths presents a tree structure. Routing protocol in which the sink has four children. Each of them leads a sub tree defined by the dotted lines. Data gathering and reconstruction of CDG are performed on the sub tree basis.

The main contributions of this propose system are:

- The spatial correlation property of a sensor node leads to inherent data sparsity in some areas, such as wavelet domain and DCT domain.
- In order to solve the sparsity of such signals, compressive sensing (CS) is exploited as a novel signal-processing paradigm that provides an efficient compressive method and recovers sparse or compressible signals.
- Spatial property of sensor node readings is exploited to strengthen the performance of networks, considering the topology structure and sensor nodes' distance.

- The spatial correlations of sensor node readings to further promote the efficiency of the data-gathering algorithm.

A. Network Model

WSNs where N sensor nodes are randomly deployed in a square area. The system model is represented by a connected graph $G(V, E)$, where the vertex set V denotes the nodes in the networks, and the edge set E denotes the wireless links between the different nodes[8]. Node i can communicate with node j if they are involved in the communication range. It assumes that the single hop distance d_{ij} between node i and node j can be represented as a Euclidean distance. At a sampling instant, each sensor node i takes a measurement x_i ; the goal of the data gathering in WSNs is to collect sufficient information to reconstruct the N-dimensional signal $X = [x_1 \dots \dots x_N]^T$. When the distance between transmission node i and receive node, j is greater than d_0 , the multi-path fading model is utilized. When the distance is less than d_0 , the free-space model is adopted.

$$E_{T_i}(L, d) = \begin{cases} E_{elec} \times L + E_{amp} \times L \times d^4, & d \geq d_0 \\ E_{elec} \times L + E_{fs} \times L \times d^2 & d < d_0 \end{cases} \quad (5)$$

$$E_{R_j}(L) = E_{elec} \times L \quad (6)$$

where $E_{T_i}(L, d)$ and $E_{R_j}(L)$ describe the energy consumption of transporting and receiving the L bit data packet. E_{elec} denotes the power expended to run the transmitter or receiver circuitry of the sensor node[9]. E_{amp} and E_{fs} represent energy consumption for a multi-path fading amplifier and free-space amplifier, respectively.

B. Modified Diffusion Wavelet

To make full use of the spatial correlation property, it takes diffusion wavelets as the sparse basis considering the spatial correlation of sensor node readings in WSNs. One is the nodes degree, and the other is the distance between the different sensor nodes. In addition, an improved QR decomposition of Givens transform is introduced to set up the sparse basis[10]. To construct the modified diffusion wavelets in detail. However, diffusion wavelets are affected significantly by the diffusion operator, which is equivalent to the wavelet function of a discrete wavelet transform. Diffusion is utilized as a smoothing and scaling technique to enable multi-scale and coarse-grained application.

Step 1: Suppose that $G(V, E)$ denotes a graph with N sensor nodes deployed in the monitoring environment. Diffusion wavelets are introduced to set up an orthonormal basis for functions supported by the topology graph of WSNs. It takes a random deployment of WSNs to explain this process.

Step 2: Calculate the weight adjacency matrix of $G(V, E)$, which is denoted as $\Omega = [w_{i,j}]$. $w_{i,j}$ is the weight of the edge in the graph[11]. Here consider two different cases of weight. The sensor node degree is chosen as the weight in the first scheme, while the is taken into consideration to exploit the spatial correlation features, aiming to mitigate the load of WSNs in another scheme. In the former case, a graph and corresponding weight adjacency matrix.

$$w_{ij} = \begin{cases} d_{ij}^Y, & i \neq j, d_{ij} \leq r \\ \chi, & otherwise \end{cases} \quad (7)$$

where r is the maximum distance among the sensor nodes that can directly communicate by a single hop. d_{ij} is the Euclidean separation between node i and node j . χ is a negative number, while χ is a small positive number.

Step 3: Generate a normalized Laplacian matrix of $G(V, E): \Lambda = [\lambda_{ij}]$. In that Λ is the degree of correlations among different function values provided at the vertices of the graph (V, E) . In the first schedule, denote λ_{ij} using Equation (8), while the other schedule considering spatial correlation implements Equation (9). Generally speaking, an eigen value or eigenvector shows the special correlations at some scale. It needs to split the space of Λ if it decomposes the signal sampled of the $G(V, E)$ in a multi-scale.

$$\lambda_{ij} = \begin{cases} 1, & i = j \\ -\frac{w_{ij}}{\sqrt{\sum_u w_{iu} \sum_u w_{uj}}}, & otherwise \end{cases} \quad (8)$$

$$\lambda_{ij} = \begin{cases} 1 - \frac{\chi}{\sum_u d_{iu}^Y} & i = j \\ -\frac{d_{ij}^Y}{\sqrt{\sum_u d_{iu}^Y \sum_u d_{uj}^Y}} & otherwise \end{cases} \quad (9)$$

Step 4: However, the diffusion operator O stems from Λ , where O shares the same eigenvalues as Λ (less than 1). The diffusion operator is $O = I - \Lambda$ or $O = \Lambda/2$; in this propose, it choose the first expression.

Step 5: Consequently, recursively raise O to power 2, and delete the diminishing eigenvalues with a threshold. Step by step, this approach splits the space spanned by the eigenvectors. Let

the initial space of O be $x_0 = \mathbb{R}^N$, which is represented by scale space $\{x_j\}_{j \in N}$ and wavelet space $\{V_j\}_{j \in N}$. Wavelet space V_j is different between x_j and x_{j+1} . Then, it derives Equation (10):

$$x_{j+1} = x_0 \oplus V_0 \oplus V_1 \oplus \dots \oplus V_j \quad (10)$$

accomplish the modified QR decomposition, where $[O]_{x_a}^{x_b}$ indicates the column space of matrix O denoted by basis x_b at scale b, and row space is denoted by basis at scale a, $[x_b]_{x_a}$ represents basis x_b denoted on the basis x_a .

Step 6: In the end, the diffusion wavelet basis Ψ is the concatenation of the scale functions and wavelet functions.

C. Modified Ant Colony Routing

In order to decrease the whole network transmission load and prolong the network lifetime, we provide a modified ant colony routing algorithm, where to speed up the convergence rate and avoid local optimal of the algorithm, pheromone impact factor is improved[12]. Here, we select the energy consumption model. The traditional ant colony optimization algorithm selects the next hop depending on Equation (11)

$$p_{ij}^g = \begin{cases} \frac{[\tau_{ij}(t)]^\zeta [\rho_{ij}(t)]^{\zeta'}}{\sum_{v \in allowed_g} [\tau_{ik}(t)]^\zeta [\rho_{ij}(t)]^{\zeta'}} & , j \in allowed_g \\ 0 & \text{hers} \end{cases} \quad (11)$$

where $\tau_{ij}(t)$ denotes the pheromone information on edge (i, j), while $\rho_{ij}(t)$ is the heuristic information on edge (i, j). ζ and ζ' are impact factors demonstrating the importance degree of the pheromone information and heuristic information. In order to speed up the convergence rate and avoid local optimal, impact factor V is modified as in Equation (12):

$$\zeta = \mu \left(1 + e^{-10 \times \left(\frac{iter}{totiter} \right)^{10}} \right) \quad (12)$$

where μ is a small positive constant $\in (0,1)$; *iter* and *totiter* refer to current iterations and total iterations, respectively. In Equation (12), V gradually becomes smaller as the number of

iterations increases. In other words, the proportion of pheromones will diminish when the number of iterations rises.

Furthermore, to yield optimal routing by the ant colony algorithm, in this subsection, a sensor node's residual energy and path length are taken into consideration simultaneously. So, the fitness value of each routing is presented as follows:

$$fitness = \beta(E_{ave} \times E_{min}) + \sigma Len_{\vartheta}^{-iter} \quad (13)$$

where E_{ave} indicates the average residual energy, while E_{min} represents the node minimal energy of ants passing through the path. Len_{ϑ}^{-iter} denotes the reciprocal of path length for given ϑ th ant and its iterations. β and σ are $\in [0,1]$ constants, and $\beta + \sigma = 1$. Consequently, the path with the largest fitness function value is chosen as the optimal routing, thus balancing the network load and prolonging the network lifetime.

D. Compressive Data Gathering

WSNs are utilized for gathering physical signal from the real world in practical applications. Without using CS theory, which is the simplest method, a data-gathering scheme with the help of the tree topology. In order to dramatically decrease communication costs and prolong the network lifetime, the authors consider that the sink node receives only M packets instead of N packets of original data from the whole network. In the end, at the sink, CS theory is used to reconstruct the original data. For the CDG algorithm, each node in the WSN multiplies its readings x_j using the corresponding j column vector of basis matrix

Φ . Next, the sensor node adds them to its own readings after receiving all same-size vectors from descendent nodes and transmitting the final results to its parent node with M packets[13]. Let us illustrate the product of CDG, where F is $M \times N$ matrix, and each column corresponds to one weight sum. In the plain CS, all nodes in WSNs transmit M packets and each has equal transmission costs; therefore, each CS measurement cost remains relatively high. An example of the plain CS mechanism. It is obvious that for these approaches (non-CS and plain CS), the former transmits fewer packets compared with plain CS from the point of view of child nodes. Provides the hybrid CS method, where non-CS is chosen when the number of packets is less than or equal to M; alternatively, plain CS is used.

According to the analysis, the network load in hybrid CS is unbalanced. Specifically, sensor nodes near the sink node will consume more energy than those far from the sink node because of forwarding data more times. This results in sensor nodes near the sink dying earlier. However, the total network costs for each random projection. To avoid the drawbacks, one can leverage the advantages of the algorithms; in this section, we present our data-gathering strategy combining joint routing and CS.

Firstly, randomly choose M projection nodes in the network with probability $\frac{M}{N}$, which follows. In the CS theory, the sink node needs M measurements to reconstruct the original data. Therefore, these M projection nodes will be selected as the gathering node, defined as g_1, g_2, \dots, g_m , to collect one random measurement y_i , and transmit y_i to the sink node. Then, distribute non-zero elements in each row of measurement matrix Φ as uniformly as possible to guarantee the sparse features of the measurement matrix; the number of non-zero elements in each row should equal to $\lfloor \frac{M}{N} \rfloor$, which is related to Algorithm 3's step 1. Additionally, each column of measurement matrix represents a sensor node, so if a column of the matrix has full zero elements, the data from its special sensor node should be thrown away. ϕ_i , the column vector of measurement matrix Φ is required to store each sensor node memory in advance.

Based on the MST algorithm, access all candidate sensor nodes of a given projection node. In the first stage, the projection node is considered one root node tree[14]. In the step 4 initialization stage in Algorithm 3, $Tree_i$ is assigned by i , and the temporary variable temp also yields i . Then steps 5–12 use the MST algorithm to construct the tree, adding the candidate nodes step by step. If temp is not empty, step 6 deletes the top node of the temp queue and puts its neighbor node in the Tree and temp. The next step is to delete them from can_i if they belong to can_i . Note that these candidate nodes must be directly connected to the parent node by a single hop[15]. If there are still some candidate nodes not involved in the tree, the Dijkstra algorithm is proposed, aiming to find the shortest path from the residual nodes to the tree (steps 14–19), and we add the residual candidate nodes.

Finally, this loop of 13–26 lines will repeat until can_i is empty. The modified ant colony routing technique is utilized to transmit packets of projection nodes to the sink node, namely step 27 of Algorithm 3. Consequently, Algorithm 3 terminates by generating the optimal routing between the projection nodes and the sink node, and an M routing tree from the projection nodes to their own candidate nodes[16]. Our novel algorithm (Algorithm 3) is shown in more detail. The modified ant colony algorithm jointly considers the sensor node's residual energy and the path length, which will not only balance the whole network load, avoiding nodes near the sink node dying earlier, but will prolong the network lifetime. In this way, the transmission costs should be greatly decreased compared to hybrid CS.

E. Algorithm Implementation

Algorithm 1: Modified diffusion wavelets.

Input: the number of sensor nodes N , communication radius r , decomposition level η , precision ϵ and MQR function.

Output: sparse basis Ψ .

1 generate a graph $G(V, E)$

2 compute weight adjacency matrix $\Omega = [w_{ij}]$ according to the vertex degree/Equation (7)

3 calculate normalized Laplacian matrix L relying on Equation (8)/Equation (9)

4 generate diffusion operator $O = I - \Lambda$

5 recursively raising O to power 2

5.1 for $\eta = 0$ to $\eta - 1$

5.2 $[x_{j+1}]_{x_j} [O]_{x_0}^{x_1} \leftarrow \text{MQR} \left([O^{2^\eta}]_{x_j}^{x_j}, \epsilon \right)$

5.3 $O_{j+1} := [O^{2^{\eta+1}}]_{x_{j+1}}^{x_{j+1}} \leftarrow [x_{j+1}]_{x_j} [O^{2^{\eta+1}}]_{x_j}^{x_j} [x_{j+1}]_{x_j}^*$

5.4 $[\Psi_\eta]_{x_j} \leftarrow \text{MQR} \left(I_{(x_j)} - [x_{j+1}]_{x_j} [x_{j+1}]_{x_j}^*, \epsilon \right)$

5.5 end for

6 concatenation of the scale functions and wavelet functions is regarded as the sparse basis Y .

MQR Function: $Q, R \leftarrow \text{MQR}(B, \epsilon)$

Input: $B: N \times N$ sparse matrix, ϵ

Output: Q, R matrix, possibly sparse, such that $B =_\epsilon QR$

(1) Q is orthogonal

(2) R is upper triangular up to a permutation

(3) The columns of Q ϵ -span the space spanned by the columns of B

Algorithm 2: Modified ant colony algorithm.

Input: the number of sensor nodes N , the power expended to run the transmitter or receiver circuitry of sensor node E_{elec} , energy consumption of multi-path fading amplifier E_{amp} , energy

consumption of free-space amplifier E_{fs} , distance threshold d_0 , impact factors of pheromone

information ς , impact factors of heuristic information ς, μ is a small positive constant $\in (0,1]$,

pheromone information on edge (i,j) τ , heuristic information on edge $(i,j)_{p,\beta}$ and σ are

$\in [0,1]$ constants [17].

Output: optimal routing Routing .

1 Initialization routing $R(\text{Path}, \text{iter}, \theta)$, energy for each node and tabu

- 2 calculate distance d_{ij} of different nodes, $\rho_{ij} = 1/d_{ij}$
- 3 while maximum iterations have not been reached
- 4 for $\theta = 1:\Theta'$
- 5 computes $allowed_s$ according to the node communication radius.
- 6 generate transition probability p_{ij}^s based on Equations (11) and (12)
- 7 choose the next hop node, relying on p_{ij}^s , modify routing and $tabu$
- 8 the destination node or not? If not, go back to step 2, or proceed to step 9
- 9 update the node residual energy based on Equations (5) and (6), routing depending on Equation (13)
- 10 end for
- 11 end while
- 12 return the optimal routing $Routing$.

Algorithm 3: Proposed algorithm.

Input: $G(V, E)$

Output: $Tree\{1, 2, \dots, M\}$ Optr

- 1 randomly select M sensor nodes s_1, s_2, \dots, s_M in the network probability $\frac{M}{N}$, generate F
- 2 for $i = 1:M$
- 3 query candidate nodes $Can_i (\phi_{ij} \neq 0)$ of projection nodes i
- 4 initializations $Tree_i \leftarrow i, temp_i \leftarrow i$
- 5 while !empty(temp) do
- 6 $CanNode \leftarrow Del(temp)$
- 7 if $CanNode$ is i 's candidate node
- 8 $Tree \leftarrow CanNode$
- 9 $temp \leftarrow CanNode$
- 10 $Del(Int_i, CanNode)$
- 11 end if
- 12 end while
- 13 while !empty(Can_i) do
- 14 for all residual candidate nodes $r \in Can_i$

```

15 Path(r) ← find a shortest path to Treei using the Dijkstra algorithm
16 if shortestpath > Path(r)
17 Shortestpath ← Path(r)
18 end if
19 end for
20 Tree ← shortestpath
21 temp ← shortestpath
22 Del(Cani, shortestpath)
23 while !empty(temp) do
24 go back to steps 7–11
25 end while
26 end while
27 optr ← Optimal routing from i to the sink node using Algorithm 2
28 return Tree{1,2 ... M}
29 end for

```

Algorithm 4: Sensor signal reconstruction.

```

1 Input: received data X, measurement matrix  $\Phi$ , the number of atom is  $n_{atom}$ 
2 Output: reconstruct data  $\hat{x}$ 
3 generate sparse basis  $\Psi$  using Algorithm 1
4 collect data  $y$  in the network using Algorithm 3
5  $A \leftarrow \Psi * \Phi$ 
6 initialization residual error  $r_0 = y, \Lambda_0 = \emptyset, A_0 = \emptyset, it = 1$ 
7 computes  $\langle r_{i-1}, a_j' \rangle$ , select the largest  $la$  values from  $\langle r_{i-1}, a_j' \rangle$ ; these values correspond to  $A'$ 's
column indexes  $j$ , constructing set  $J_0$ 
8 set  $\Lambda_{it} = \Lambda_{it-1} \cup J_0, A_{it} = A_{it-1} \cup a_j$  (for all  $j \in J_0$ )
9  $\hat{x}_{it} = (A_{it}^T A_{it})^{-1} A_{it}^T y$ 
10 updates  $r_{it} = y - A_{it} \hat{x}_{it}$ 
11  $it = it + 1$ , if  $it \leq K$  go back to step 7, or proceed to step 12

```

12 reconstruct \hat{x} , which is the generation value of the last iteration \hat{x}_{i_T} .

IV. RESULT AND DISCUSSION

To evaluate the performance of our scheme by experiments. Ii evaluate our scheme mainly in terms of the sparse basis comparison; the reconstruction performance of the novel mechanism; the reconstruction error for different schemes; the energy consumption based on non-CS, plain CS, hybrid CS and our proposed algorithm (sparse basis is based on distance); and network lifetime performance between the different schemes and our algorithms[19]. In our simulations, all programs have been run in the NS2 platform. Moreover, $E_{elec} = 50 \text{ nJ/bit}$, $E_{amp} = 60 \text{ pJ/bit/m}^4$, $E_{fs} = 100 \text{ pJ/bit/m}^2$, $L = 1024 \text{ bits}$, initial energy $E_0 = 5j$.

Table 1. Type Styles

Table 1. Simulation Parameters

Parameter name	Parameter value
Stimulation tool	NS2
Packet size	512kb
Channel	Wireless
Number of Mobile nodes	80
Communication agent	UDP
MAC type	802-11

NS2: The Network Simulator (ns2) is a discrete event driven simulator developed at UC Berkeley. We are using Network Simulator NS2 for simulations of protocols[18]. It provides considerable support for simulation of TCP, routing and multicast protocols over wired and wireless networks. Ns-2 code is written either in C++ and OTCL and is kept in a separate file that is executed by OTCL interpreter, thus generating an output file for NAM (Network animator). It then plots the nodes in a position defined by the code script and exhibits the output of the nodes communicating with each other.

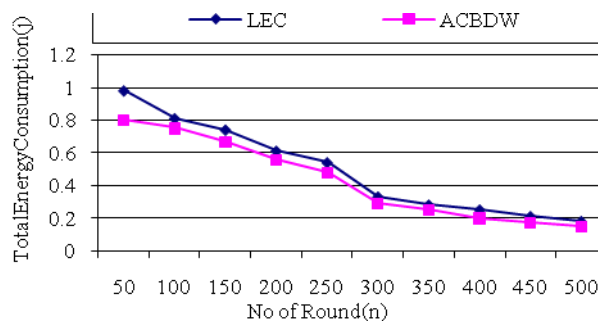


Fig1. Comparison of no of round vs total energy consumption

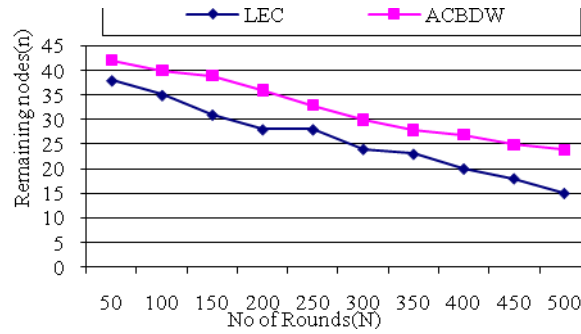


Fig2. Comparison of no of round vs remaining nodes

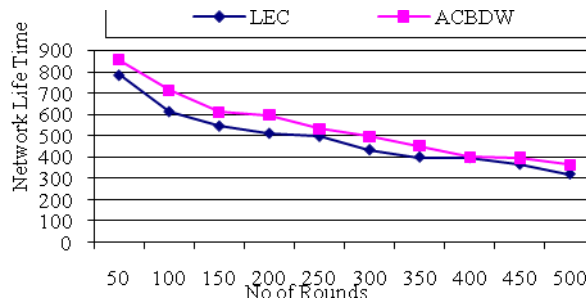


Fig3. Comparison of no of round vs network life time

V. CONCLUSIONS

Therefore, in this mechanism, diffusion wavelets based on sensor nodes’ degree and different nodes’ distance considering the above factors are proposed. Additionally, to further reduce the transport costs in WSNs, a sparse measurement matrix is utilized and MST and modified ant colony routing are jointly applied to mitigate energy consumption and balance the network load, especially lowering the transmission costs for those nodes nearest the sink node. Experimental results have shown that our sparse basis can sparsity the signal well. This method can also accurately reconstruct the original signal. Moreover, the reconstruction error of our scheme is less than DFT [20]. Compared with existing data-gathering approaches, our proposed algorithm not only minimizes the energy consumption of the network, but prolongs the network lifetime.

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